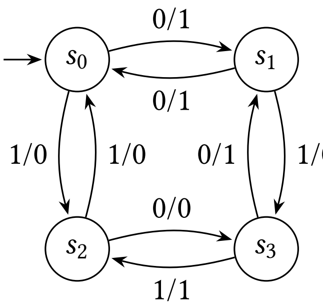
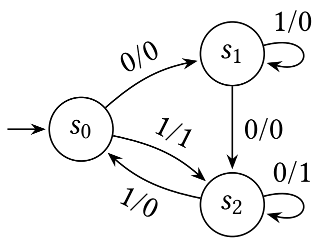
1. For each given regular expression P, construct a DFA (Deterministic Finite Automaton), and find the number of accepted word of length at most 5, i.e. the size of the set $\mathcal{L}' = \{w \in \mathcal{L}(P) \mid |w| \leq 5\}$. For "any" (`.`) and "negative" (`[^.]`) matches, assume that the alphabet is $\Sigma = \{\mathtt{a}, \mathtt{b}, \mathtt{c}, \mathtt{d}\}$.
   (a) $P_1 =$ `ab*`
   (b) $P_2 =$ `a+b?c`
   (c) $P_3 =$ `[^cd]+c{3}`
   (d) $P_4 =$ `[^a](.|ddd)?`
   (e) $P_5 =$ `d(a|bc)*`
   (f) $P_6 =$ `((a|ab)[cd]){2}`

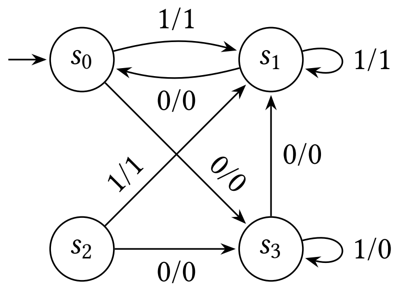
2. Describe the set of strings defined by each of these sets of productions in EBNF (extended Backus-Naur form).
   (a) $\langle \mathit{string}\rangle ::= \langle L\rangle+\ \langle D\rangle?\ \langle L\rangle+$
   $\langle L\rangle ::= \mathtt{a} \mid \mathtt{b} \mid \mathtt{c}$
   $\langle D\rangle ::= \mathtt{0} \mid \mathtt{1}$
   (b) $\langle \mathit{string}\rangle ::= \langle \mathit{sign}\rangle?\ \langle N\rangle$
   $\langle \mathit{sign}\rangle ::= \text{'+'} \mid \text{'–'}$
   $\langle N\rangle ::= \langle D\rangle\ (\langle D\rangle \mid \mathtt{0})^*$
   $\langle D\rangle ::= \mathtt{1} \mid \ldots \mid \mathtt{9}$
   (c) $\langle \mathit{string}\rangle ::= \langle L\rangle^*\ (\langle D\rangle+)?\ \langle L\rangle^*$
   $\langle L\rangle ::= \mathtt{x} \mid \mathtt{y}$
   $\langle D\rangle ::= \mathtt{0} \mid \mathtt{1}$
   (d) $\langle \mathit{string}\rangle ::= \langle C\rangle\ \langle R\rangle^*$
   $\langle C\rangle ::= \mathtt{a} \mid \ldots \mid \mathtt{z} \mid \mathtt{A} \mid \ldots \mid \mathtt{Z}$
   $\langle D\rangle ::= \mathtt{0} \mid \ldots \mid \mathtt{9}$
   $\langle R\rangle ::= \langle C\rangle \mid \langle D\rangle \mid \text{'\_'}$

3. Let $\mathcal{G} = \langle V, T, S, P\rangle$ be the phrase-structure grammar with vocabulary $V = \{\mathtt{A}, \mathtt{S}\}$, terminal symbols $T = \{\mathtt{0}, \mathtt{1}\}$, start symbol $S = \mathtt{S}$, and set of productions P: $\mathtt{S} \longrightarrow \mathtt{1S}$, $\mathtt{S} \longrightarrow \mathtt{00A}$, $\mathtt{A} \longrightarrow \mathtt{0A}$, $\mathtt{A} \longrightarrow \mathtt{0}$.
   (a) Show that `111000` belongs to the language generated by $\mathcal{G}$.
   (b) Show that `11001` does not belong to the language generated by $\mathcal{G}$.
   (c) What is the language generated by $\mathcal{G}$?

4. Find the output generated from the input string `01110` for each of the following Mealy machines.







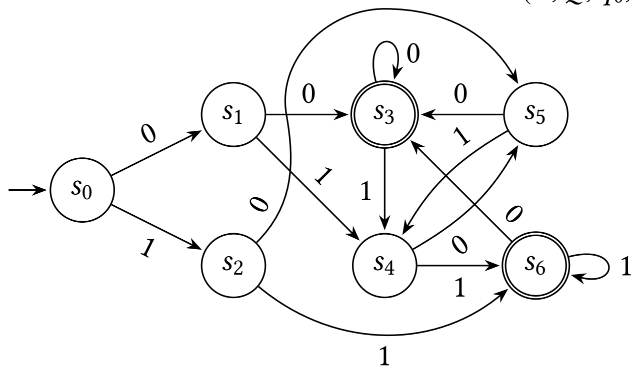
5. Construct a Moore machine for each of the following descriptions.
   (a) Determine the residue modulo 3 of the input treated as a binary number. For example, for input ε (which corresponds to "value" 0) the residue is 0; `101` (5 in decimal) has residue 2; and `1010` (value 10) has residue 1.
   (b) Output the residue modulo 5 of the input from $\{0, 1, 2\}^*$ treated as a ternary (base 3) number.
   (c) Output A if the binary input ends with `101`; output B if it ends with `110`; otherwise output C.

6. Show that regular languages are *closed* under the following operations.
   (a) Union, that is, if L_1 and L_2 are regular languages, then $L_1 \cup L_2$ is also regular.
   (b) Concatenation, that is, if L_1 and L_2 are regular languages, then $L_1 \cdot L_2$ is also regular.
   (c) Kleene star, that is, if L is a regular language, then L^* is also regular.
   (d) Complement, that is, if L is a regular language, then $\overline{L} = \Sigma^* - L$ is also regular.
   (e) Intersection, that is, if L_1 and L_2 are regular languages, then $L_1 \cap L_2$ is also regular.

7. Determine whether the following languages are regular or not. For non-regular languages, use Pumping lemma to prove that they are not regular. For each regular language, provide a regular expression and construct an ε-NFA.
   (a) $L_1 = \{w \in \{0,1\}^* \mid \text{length of } w \text{ is odd}\}$
   (b) $L_2 = \{0^n 1^n \mid n \in \mathbb{N}\}$
   (c) $L_3 = \{w \in \{0,1\}^* \mid w \text{ contains an even number of 1s}\}$
   (d) $L_4 = \{1^{n^2} \mid n \in \mathbb{N}\}$

8. Consider a finite-state automaton $M = (\Sigma, Q, q_0, F, \delta)$ and a non-negative integer k. Let R_k be the relation on the set of states of M such that $s\ R_k\ t$ if and only if for every input string $w \in \Sigma^*$ with $|w| \le k$, $\delta(s,w)$ and $\delta(t,w)$ are both final states or both not final states. Furthermore, let R^* be the relation on the set of states of M such that $s\ R^*\ t$ if and only if for every input string $w \in \Sigma^*$, regardless of length, $\delta(s,w)$ and $\delta(t,w)$ are both final states or both not final states.
   (a) Show that for every nonnegative integer k, R_k is an equivalence relation on S.
   Two states s and t are called k-equivalent if $s\ R_k\ t$.
   (b) Show that R^* is an equivalence relation on S.
   Two states s and t are called $*$-equivalent if $s\ R^*\ t$.
   (c) Show that if two states s and t are k-equivalent ($k > 0$), then they are also $(k-1)$-equivalent.
   (d) Show that the equivalence classes of R_k are a *refinement* of the equivalence classes of R_{k-1}.
   (e) Show that if two states s and t are k-equivalent for every non-negative integer k, then they are $*$-equivalent.
   (f) Show that all states in a given R^*-equivalence class are final or all are not final.
   (g) Show that if two states s and t are $*$-equivalent, then $\delta(s,a)$ and $\delta(t,a)$ are also $*$-equivalent for all $a \in \Sigma$.

9. Consider the finite-state automaton $M = (\Sigma, Q, q_0, F, \delta)$ depicted below.



   (a) Find the k-equivalence classes of M for $k = 0, 1, 2, 3$.
   (b) Find the $*$-equivalence classes of M.
   (c) Construct the quotient automaton $\overline{M}$ of M.
   ▹ The quotient automaton $\overline{M}$ of the deterministic finite-state automaton $M = (\Sigma, S, s_0, F, \delta)$ is the finite state automaton $\overline{M} = (\Sigma, \overline{S}, [s_0]_{R^*}, \overline{F}, \overline{\delta})$, where the set of states $\overline{S}$ is the set of R^*-equivalence classes of S; the transition function $\overline{\delta}$ is defined by $\overline{\delta}([s]_{R^*}, a) = [\delta(s,a)]_{R^*}$ for all states $[s]_{R^*}$ of $\overline{M}$ and input symbols $a \in \Sigma$; and $\overline{F}$ is the set consiting of R^*-equivalence classes of final states of M.

10. Solve the following regex crosswords[1]. Fill each cell with a single ASCII character (an uppercase letter, a digit, a punctuation mark, or a space). Each row/column, when read left to right or top to bottom must match the regular expression(s) given for that row/column.

	[^SPEAK]+	EP|IP|EF
HE|LL|O+		
[PLEASE]+		

	(FI|A)+	(YE|OT)K	(.)[IF]+	[NODE]+	(FY|F|RG)+
(Y|F)(.)\2[DAF]\1					
(U|O|I)*T[FRO]+					
[KANE]*[GIN]*					

	[^NRU](NO|ON)	(D|FU|UF)+	(FO|A|R)*	(N|A)*
[RUNT]*				
O.*[HAT]				
(.)*DO\1				

	[^MCI]+	.A	(TM|BF)	
(CAT|A-T)+				[^KI\sP]+
[MA\-\sE]+				(M|APS|EA)*
	[AI][E\s]	[A\-Z]+	[\sT\-M]+	

	(.)\1(.)\2	[C\sOU]+	[^PU\sH]+	
.[LUH]+				.*L+
(P|K)[^U]+				[PUF\s]*
.*C+[TIF]				[TIC]*
(NO|ONE|ION)*				[NOI\sE]+
	[PIF]+	.*[OWE]*	(TN|LF|TF)*	

	.*E.*	[^P]I(IT|ME)	(EM|FE)(IT|IP)	(TS|PE|KE)*
(EP|ST)*				
T[A-Z]*				
.M.T				
.*P.[S-X]+				

[1] Credits: https://regexcrossword.com